Markov-chain Monte Carlo: A modern primer Lecture 1: Fundamentals Part 2/2: Single particle on the path graph

Werner Krauth Département de physique, Ecole normale supérieure Paris, France

A Set of Lectures University of Kent Canterbury, Great Britain

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Metropolis algorithm on path graph (1/5)



- Path graph \mathcal{P}_n so that $\Omega_n = \{1, \ldots, n\}$.
- Phantom vertices and edges.

Metropolis algorithm (NB: $P_{ij} = A_{ij}P_{ij}$ for $i \neq j$):

1 Move set
$$\mathcal{L} = \{+, -\}$$
.

2
$$\mathcal{A}$$
 flat $\rightarrow \sigma = \text{choice}(\mathcal{L}).$

Solution Metropolis filter: Accept with probability min $(1, \pi_i/\pi_i)$.

Metropolis algorithm on path graph (2/5)



Detailed balance:

$$\underbrace{\pi_i \mathbf{P}_{ij}}_{\mathcal{F}_{ij}} = \underbrace{\pi_j \mathbf{P}_{ji}}_{\mathcal{F}_{ji}}$$

Metropolis algorithm:

$$\mathcal{F}_{ij} = \frac{1}{2} \min \left(\pi_i, \pi_j \right) \Leftrightarrow \mathcal{P}_{ij} = \frac{1}{2} \min \left(1, \pi_j / \pi_i \right)$$

• Metropolis filter (NB: $P_{ij} = A_{ij}P_{ij}$):

$$\mathcal{P}_{ij} = \min\left(\mathbf{1}, \pi_j/\pi_i\right)$$

Metropolis algorithm on path graph (3/5)



• Global balance ($\pi_i = \sum_j \pi_j P_{ji} = \sum_j \mathcal{F}_{ji}$):



- Irreducibility OK if no holes in π .
- Aperiodicity OK, thanks to boundaries

Metropolis algorithm on path graph (4/5)



procedure metro-path input x $\sigma \leftarrow \text{choice}(\mathcal{L}) \ (\mathcal{L} = \{-1, +1\})$ if $(\operatorname{ran}(0, 1) < \pi_{x+\sigma}/\pi_x)$ then $\begin{cases} x_i \leftarrow x_i + \sigma \\ \text{output } x \end{cases}$

Metropolis algorithm on path graph (5/5)



Model	Definition	Conductance	t _{mix}
Flat	$\pi_i = \frac{1}{n}$	$\mathcal{O}(1/n)$	$\mathcal{O}\left(n^{2}\right)$
Square	$\pi_{2k-1} = \frac{2}{3n}, \ \pi_{2k} = \frac{4}{3n}$	\sim 2/(3 n)	$\mathcal{O}(n^2)$
V-shape	$\pi_i = \frac{4}{n^2} \left \frac{n+1}{2} - i \right ^{-1}$	\sim 2 $/n^2$	$\mathcal{O}\left(n^2\log n\right)$

NB: Graph diameter *n*. NNB: π normalized. NNNB: Bottleneck between $i = \frac{n}{2}$ and $j = \frac{n}{2} + 1$.

Lifting on the path graph (1/7)

Probability distribution $\pi = (1/n, ..., 1/n)$ (Diaconis et al. 2000)

• "Collapsed" Markov chain:



• "Lifted" Markov chain $\hat{\Omega} = \Omega \times \{-,+\}$:



• Irreducible, but not aperiodic. $\hat{\pi}_{i,\sigma} = \frac{1}{2}\pi_i = 1/(2n)$

Lifting on the path graph (2/7)

Probability distribution $\pi = (1/n, ..., 1/n)$ (Diaconis et al. 2000)

• Transport + resampling



• "Lifted" Markov chain $\hat{\Omega} = \Omega \times \{-,+\}$:



Lifting on the path graph (3/7)



- Diaconis, Holmes, Neal (2000).
- Ω (samples) + \mathcal{L} (moves) $\rightarrow \hat{\Omega} = \Omega \times \mathcal{L}$ (lifted samples)
- Resampling with rate 1/n.
- Phantoms illustrate the "rejection \rightarrow lifting" mystery.
- Let's generalize.

Lifting on the path graph (4/7)

General probability distribution $\pi = (\pi_1, \ldots, \pi_n)$

• "Collapsed" Markov chain:



• "Lifted" Markov chain $\hat{\Omega} = \Omega \times \{-,+\}$:



• Replace all rejections by lifting moves.

Lifting on the path graph (5/7)

• "Lifted Markov chain: Transport"

NB: The
$$\frac{1}{2} \Leftrightarrow \hat{\pi}_{i,\sigma} = \frac{1}{2}\pi_i$$

• "Lifted Markov chain: Resampling"

$$(\mathbf{i},+1)$$

$$\frac{1}{2}\pi_i \epsilon \downarrow \uparrow \frac{1}{2}\pi_i \epsilon$$

$$(\mathbf{i},-1)$$

• Resampling can often be dropped

Lifting on the path graph (6/7)

• Transport

procedure transport-path input $\{x, \sigma\}$ (configuration $\in \widehat{\Omega} = \Omega \times \{+, -\}$) if $(\operatorname{ran}(0, 1) < \pi_{x+\sigma}/\pi_x)$ then $\{x_i \leftarrow x_i + \sigma$ else $\{\sigma \leftarrow -\sigma$ output $\{x, \sigma\}$

Resampling

procedure resample-path input $\{x, \sigma\}$ (configuration $\in \widehat{\Omega} = \Omega \times \{+, -\}$) if $(\operatorname{ran}(0, 1) < p$) then (p: resampling rate) $\{\sigma \leftarrow -\sigma$ output $\{x, \sigma\}$

• Mix freely!

Lifting on the path graph (7/7)



Model	Conductance	t _{mix} (collapsed)	t _{mix} (lifted)
Flat	<i>O</i> (1/ <i>n</i>)	$\mathcal{O}\left(n^{2}\right)$	$\mathcal{O}(n)$
Square	$\mathcal{O}(1/n)$	$\mathcal{O}(n^2)$	$\mathcal{O}\left(n^{2}\right)$
V-shape	$\mathcal{O}\left(1/n^2\right)$	$\mathcal{O}\left(n^2 \log n\right)$	$\mathcal{O}\left(n^{2}\right)$

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